

## ROOT SQUARE MEAN LABELING OF ARROW GRAPHS WITH ENCODING AND DECODING

T. CHRISTY<sup>1\*</sup>, G. PALANI<sup>2</sup>, §

**ABSTRACT.** A graph  $G$  with  $p$  vertices and  $q$  edges is called a Root Square mean labeling if it is possible to label the vertices  $x \in v$  with distinct labels  $\rho(x)$  from  $1, 2, \dots, q+1$  in such a way that each edge  $e = ab$  is labeled with  $\rho = \left\lceil \sqrt{\frac{\rho(a)^2 + \rho(b)^2}{2}} \right\rceil$  or  $\left\lfloor \sqrt{\frac{\rho(a)^2 + \rho(b)^2}{2}} \right\rfloor$  then the edge labels are distinct. In this case  $\rho$  is called Root Square mean (*RSM*) labeling of  $G$ . In this paper we prove Arrow graphs  $A_n^2, A_n^3, A_n^4$  admits Root Square mean (*RSM*) labeling. In today's world, digital data transfer is becoming more and more common in all industries. Data security plays a critical role in the delivery and storage of data. Labeling is an essential component of the cryptosystem. A new approach in the encoding and decoding process on the Root Square Mean Labeling is applied in this paper through an algorithm for encoding and decoding of secured message.

**Keywords:** Root Square mean (*RSM*) labeling, Arrow graphs

**AMS Subject Classification:** 05078

### 1. INTRODUCTION

Consider  $(p, q)$  graph  $G = (V, E)$ , where the vertex set and edge set of the graph  $G$  are denoted, respectively, by the symbols  $V(G)$  and  $E(G)$ . For the basic terminology and notations we refer to Harary [2]. Rosa proposed the concept of graph labeling in 1967. Gallian [1] offers a thorough analysis of graph labeling. The concept of Mean labeling of graph was introduced by Ponraj et al. [3]. The root square mean labeling was introduced by Sandhya et al. and they have proved Path, Cycle, comb, Ladder, Triangular snake, Quadrilateral Snake, Complete graph [6, 7]. Further Meena et al. investigated in their paper studied on some cycle related graphs and theta graphs and proved that the graphs admit Root Mean Square labeling [4, 5]. Cryptography is a technique for information protection that converts an original communication into a coded text that can be decoded by a knowledgeable person. We require a secret key in order to convert a regular letter into a secret encrypted

<sup>1</sup> Department of Mathematics, Patrician College of Arts and Science, Affiliated to University of Madras, Chennai, India.

e-mail: debujone@gmail.com; ORCID: <https://orcid.org/0009-0009-9797-2490>.

<sup>2</sup> Department of Mathematics, Dr. Ambedkar Government Arts College, Chennai, India.

e-mail: gpalani32@yahoo.co.in; ORCID: <https://orcid.org/0000-0003-0909-8254>.

\* Corresponding author.

§ Manuscript received: April 08, 2024; accepted: July 01, 2024; TWMS Journal of Applied and Engineering Mathematics, Vol.15, No.6; © Işık University, Department of Mathematics, 2025; all rights reserved.

one. Saverio Caminiti first proposed the concept of coding languages[8]. The encoding and decoding algorithm using AUM Block sum labeling for path and other graphs were discussed by Uma Maheswari et al. [9]

## 2. PRELIMINARIES

**Definition 2.1** (Mean Labeling). A function  $\rho$  is called mean labeling for a graph  $G=(V, E)$  if  $\rho : V \rightarrow \{0, 1, 2, 3, \dots, q\}$  is injective and the induce function  $\rho^* : E \rightarrow \{1, 2, 3, \dots, q\}$  defined as  $\rho^* = \left\lceil \frac{\rho(a)+\rho(b)}{2} \right\rceil$  or  $\left\lfloor \frac{\rho(a)+\rho(b)}{2} \right\rfloor$  is bijective for every edge. A graph  $G$  is called mean labeling.

**Definition 2.2** (Root Square Mean Labeling). A graph  $G$  with  $p$  vertices and  $q$  edges is called a RSM labeling. if it is possible to label the vertices  $x \in v$  with distinct labels  $\rho(x)$  from  $1, 2, \dots, q+1$  in such a way that each edge  $e = ab$  is labeled with  $\rho = \left\lceil \sqrt{\frac{\rho(a)^2+\rho(b)^2}{2}} \right\rceil$  or  $\left\lfloor \sqrt{\frac{\rho(a)^2+\rho(b)^2}{2}} \right\rfloor$  then the edge labels are distinct. In this case  $\rho$  is called Root Square mean(RSM) labeling of  $G$

**Definition 2.3** (Arrow graph). An arrow graph  $A_n^t$  with width  $t$  and length  $n$  is obtained by joining a vertex  $v$  with superior vertices of  $P_m \times P_n$  by  $m$  new edges from one end.

## 3. MAIN RESULTS

**Theorem 3.1.** Arrow Graph  $A_n^2$  is a RSM labeling.

*Proof.* Let  $G$  be an arrow graph  $A_n^2$ .

Let  $V(G) = \{a_i, b_i, c; 1 \leq i \leq n\}$  and  $E(G) = \{(a_i a_{i+1}); (b_i b_{i+1}); 1 \leq i \leq n-1\} \cup \{a_i b_i; 1 \leq i \leq n\} \cup \{(a_1 c), (b_1 c)\}$

Define a function  $\rho : V \rightarrow \{1, 2, 3, \dots, q+1\}$

Let us label the vertices as follows

$\rho(a_1) = 3; \rho(a_{i+1}) = 3i + 4; 1 \leq i \leq n-1$

$\rho(b_i) = 3i - 1; \rho(c) = 1; 1 \leq i \leq n$

Then the edges are labeled with

$\rho^*(a_i a_{i+1}) = 3i + 2; \rho^*(b_i b_{i+1}) = 3i + 1; 1 \leq i \leq n-1$

$\rho^*(a_i b_i) = 3i; 1 \leq i \leq n \rho^*(cb_1) = 1; \rho^*(a_1 c) = 2$

Then we get distinct edge labels.

Hence Arrow Graph  $A_n^2$  is a RSM Labeling. □

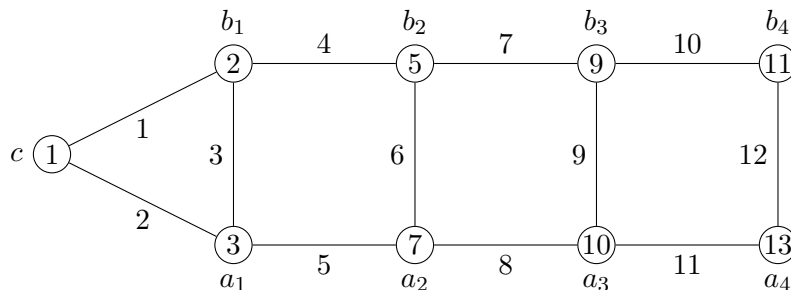


FIGURE 1. Arrow Graph  $A_3^2$

**Theorem 3.2.** Arrow Graph  $A_n^3$  is a RSM labeling.

*Proof.* Let  $G$  be an arrow graph  $A_n^3$ , Let  $V(G) = \{a_i, b_i, c_i, d; 1 \leq i \leq n\}$  and  $E(G) = \{(a_i a_{i+1}); (b_i b_{i+1}); 1 \leq i \leq n-1\} \cup \{(b_i a_i); (c_i b_i); (c_i c_{i+1}); 1 \leq i \leq n\} \cup \{(a_1 d), (c_1 d)\}$

Define a function  $\rho : V \rightarrow \{1, 2, 3, \dots, q+1\}$

Let us label the vertices as follows

$$\rho(a_1) = 3; \rho(a_{i+1}) = 5i + 5; 1 \leq i \leq n-1$$

$$\rho(b_1) = 5; \rho(b_{2i}) = 10i - 2; \rho(b_{2i+1}) = 10i + 3; 1 \leq i \leq n-1$$

$$\rho(c_1) = 2; \rho(c_{2i}) = 10i - 3; \rho(c_{2i+1}) = 10i + 1; 1 \leq i \leq n-1$$

Then the edges are labeled with

$$\rho^*(a_1 a_2) = 7; \rho^*(a_{i+1} a_{i+2}) = 5i + 8; \rho^*(b_i b_{i+1}) = 5i + 1; 1 \leq i \leq n-1$$

$$\rho^*(c_i c_{i+1}) = 5i; 1 \leq i \leq n-1$$

$$\rho^*(a_i b_i) = 5i - 1; 1 \leq i \leq n$$

$$\rho^*(b_1 c_1) = 3; \rho^*(b_{2i} c_{2i}) = 8; \rho^*(b_i c_i) = 5i + 7; 3 \leq i \leq n-2$$

$$\rho^*(a_1 d) = 1; \rho^*(c_1 d) = 2$$

Then we get distinct edge labels.

Hence Arrow Graph  $A_n^3$  is a root square mean graph.  $\square$

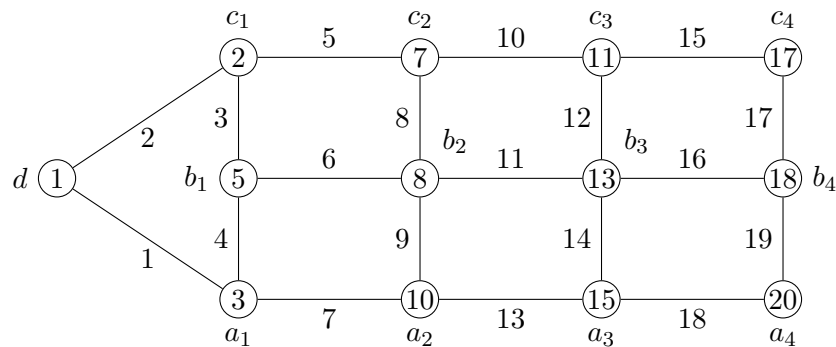


FIGURE 2. Arrow Graph  $A_3^3$

**Theorem 3.3.** Arrow Graph  $A_n^4$  is a RSM labeling.

*Proof.* Let  $G$  be an arrow graph  $A_n^4$ .

Let  $V(G) = \{a_i, b_i, c_i, d_i, e; 1 \leq i \leq n\}$  and  $E(G) = \{(ed_1, ea_1); (a_i a_{i+1}); (b_i b_{i+1}); (c_i c_{i+1}); (d_i d_{i+1}); 1 \leq i \leq n-1\} \cup \{(b_i a_i); (c_i b_i); (c_i d_i); 1 \leq i \leq n\}$

Define a function  $\rho : V \rightarrow \{1, 2, 3, \dots, q+1\}$

Let us label the vertices as follows

$$\rho(a_1) = 3; \rho(a_{i+1}) = 7i + 6; 1 \leq i \leq n-1$$

$$\rho(b_{3i-2}) = 21i + 16; \rho(b_{3i-1}) = 21i - 10; \rho(b_{3i}) = 21i - 3; 1 \leq i \leq n$$

$$\rho(c_{3i-2}) = 21i - 17; \rho(c_{3i-1}) = 21i - 11; \rho(c_{3i}) = 21i - 4; 1 \leq i \leq n$$

$$\rho(d_2) = 8; \rho(d_{2i+1}) = 14i + 1; \rho(d_{2i+2}) = 14i + 9; 1 \leq i \leq n$$

Then the edges are labeled with  $\rho^*(a_1 e) = 2; \rho^*(d_1 e) = 1$

$$\rho^*(a_i a_{i+1}) = 7i + 2; \rho^*(b_i b_{i+1}) = 7i + 1; 1 \leq i \leq n-1$$

$$\rho^*(c_i c_{i+1}) = 7i; \rho^*(d_i d_{i+1}) = 7i - 1; 1 \leq i \leq n-1$$

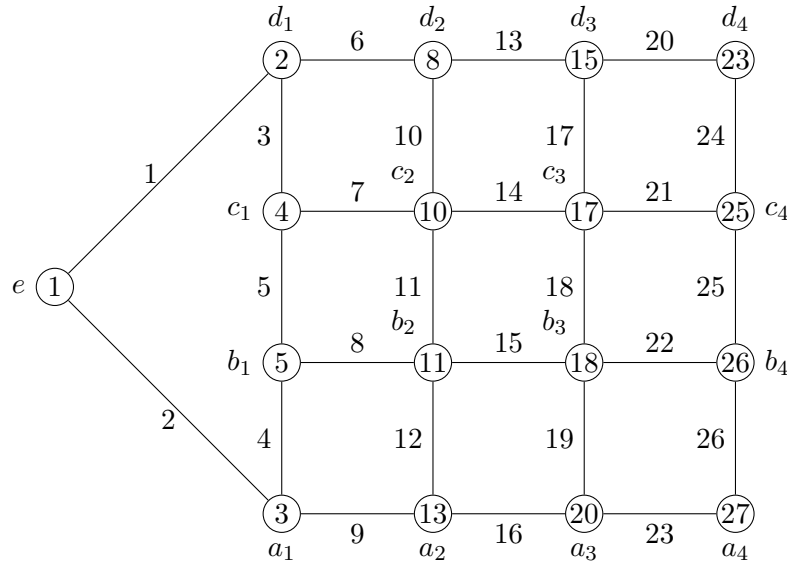
$$\rho^*(a_1 b_1) = 4; \rho^*(a_{i+1} b_{i+1}) = 7i + 5; 1 \leq i \leq n-1$$

$$\rho^*(b_1 c_1) = 5; \rho^*(b_{i+1} c_{i+1}) = 7i + 4; 1 \leq i \leq n-1$$

$$\rho^*(c_i d_i) = 7i - 4; 1 \leq i \leq n$$

Then we get distinct edge labels.

Hence Arrow Graph  $A_n^4$  is a RSM labeling.  $\square$

FIGURE 3. Arrow Graph  $A_3^4$ 

## 4. APPLICATION

**Definition 4.1** (Plain text). *An original intelligible message is called as Plain text.*

**Definition 4.2** (Cipher text). *The Transformed message after coding is called as Cipher text.*

In this section, We used Root Square Mean Labeling of Arrow Graph  $A_2^3$  to encode a message and created novel encoding and decoding techniques that increase the secrecy of the coded message.

**4.1. Algorithm for Encoding Arrow Graph. Step.1** Number the 26 alphabets in multiples of five. Five represent vowels, and the remaining letters represent consonants as follows,

The vowels a,e,i,o,u is assigned to multiples of five ie) 5,10,15,20,25 and other numbers 1,2,3,4,6,7,8,9,11,12,13,14,16,17,18,19,21,22,23,24,26 are assigned to consonants in the order.

**Step.2** Each character is given a shift cipher using the formula  $u_i = (u + n)(\text{mod } 26)$ , which moved each number n places. Here, n is the message's length and u is the number assigned to that character.

**Step.3** Find the smallest positive integer  $v_i$  such that  $w_j$  is the geometric mean of  $u_i$  and  $v_i$ , and take it as such.

**Step.4** Select the Arrow Graph  $A_2^3$  and label the edges as  $v_i$  and the remaining edges as  $w_j$  in the order.

**Step.5** Reassign the edge labels as  $E_i = v_i + E(j)$  for  $i = 1, 2, \dots, n$  and  $E_i = w_j + E(j)$  for  $j = n + 1, n + 2, \dots, 2n$  where  $E(i)$  is the edges of the  $A_2^3$ .

**Step.6** The secret to decoding the message is to deliver the cipher text as  $A_2^3$ , along with the updated edge label.

Now we find that the edge labels of Arrow graph is encoded and a new cipher text is

obtained. We proceed to decode the cipher text as follows using the labels of the edges of Key Graph .

**4.2. Algorithm for Decoding the Cipher text of Arrow Graph. Step.1** Find the Root square mean labeling for the given Key graph .

**Step.2** Determine the integers  $w_j = E_j - E(j)$  for  $j = n + 1, n + 2, \dots, 2n$  and  $v_i = E_i - E(i)$  for  $i = 1, 2, \dots, n$ , where the labels in the edges of the key graph are  $E'_i$ 's and the edges of the same graph are  $E(i)$ 's.

**Step.3** Identify the positive number. Using the relation,  $u_i = \frac{w_j^2}{v_i}$  for  $i = 1, 2, 3, 4, 5, 6$  and  $j = 7, 8, 9, 10, 11, 12$

**Step.4** Find the value of  $A_i$  using  $A_i = [u_i - \frac{n}{2}] \bmod 26$  Where  $n$  represents the number of edges and the alphabet for each  $A_i$  from the encoding table to get the plain text.

Now we illustrate the application through an example by assigning the edge labels of Arrow graph  $A_3^2$  by "CHURCH". The following is the encoding procedure.

**Encoding**

A	B	C	D	E	F	G	H	I	J	K	L	M
5	1	2	3	10	4	6	7	15	8	9	11	12
N	O	P	Q	R	S	T	U	V	W	X	Y	Z
13	20	14	16	17	18	19	25	21	22	23	24	26

TABLE 1. Encoding Table

Length of the word is 6  
Using the shift cipher  $u_i = (u + n) \pmod{26}$   
We get

C	H	U	R	C	H
2	7	25	17	2	7
8	13	5	23	8	13

Denote  $u_1 = 8, u_2 = 13, u_3 = 5, u_4 = 23, u_5 = 8, u_6 = 13$ .

For  $i = 1 \dots n$  and  $j = n + 1, n + 2, \dots, 2n$ , find the lowest positive integer  $V_i$  such that the geometric mean between  $u_i$  and  $v_i$  is a positive integer. Denote this as  $w_j$ .

$$G.M(u_1, v_1) = G.M(8, 8) = 8 = w_7; G.M(u_2, v_2) = G.M(13, 13) = 13 = w_8$$

$$G.M(u_3, v_3) = G.M(5, 5) = 5 = w_9; G.M(u_4, v_4) = G.M(23, 23) = 23 = w_{10}$$

$$G.M(u_5, v_5) = G.M(8, 8) = 8 = w_{11}; G.M(u_6, v_6) = G.M(13, 13) = 13 = w_{12}$$

Reassign the edge labels as

$E_i = v_i + E(i)$  for  $i = 1, 2, 3, 4, 5, 6$  and  $E_j = w_j + E(j)$  for  $j = 7, 8, 9, 10, 11, 12$  where  $E(i)$  is the edges label of  $An^2$ .

$$E_1 = v_1 + E(1) = 8 + 1 = 9; E_2 = v_2 + E(2) = 13 + 2 = 15$$

$$E_3 = v_3 + E(3) = 5 + 3 = 8; E_4 = v_4 + E(4) = 23 + 4 = 27$$

$$E_5 = v_5 + E(5) = 8 + 5 = 13; E_6 = v_6 + E(6) = 13 + 6 = 19$$

$$E_7 = w_7 + E(7) = 8 + 7 = 15; E_8 = w_8 + E(8) = 13 + 8 = 21$$

$$E_9 = w_9 + E(9) = 5 + 9 = 14; E_{10} = w_{10} + E(10) = 23 + 10 = 33$$

$$E_{11} = w_{11} + E(11) = 8 + 11 = 19; \quad E_{12} = w_{12} + E(12) = 13 + 12 = 25$$

The Arrow graph  $A_n^2$  with the relabeled edges is given below  
Send the receiver the labeled graph above so that they can decode it.

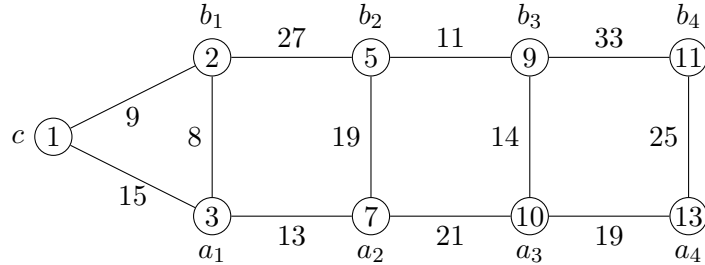


FIGURE 4. Arrow Graph  $A_3^2$

### Key for decoding

The message can be decoded using the labeled graph seen in figure 4.

### Decoding

- After receiving the labeled graph find the plain text using the above decoding algorithm.
- Finding the integers  $v_i$  and  $w_j$  using the key graph figure 4 and the relation,  $v_i = E_i - E(i)$  for  $i = 1, 2, 3, 4, 5, 6$  and  $w_j = E_j - E(j)$  for  $j = 7, 8, 9, 10, 11, 12$   
We get

$$v_1 = E_1 - E(1) = 9 - 1 = 8; \quad v_2 = E_2 - E(2) = 15 - 2 = 13$$

$$v_3 = E_3 - E(3) = 8 - 3 = 5; \quad v_4 = E_4 - E(4) = 27 - 4 = 23$$

$$v_5 = E_5 - E(5) = 13 - 5 = 8; \quad v_6 = E_6 - E(6) = 19 - 6 = 13$$

$$w_7 = E_7 - E(7) = 15 - 7 = 8; \quad w_8 = E_8 - E(8) = 21 - 8 = 13$$

$$w_9 = E_9 - E(9) = 14 - 9 = 5; \quad w_{10} = E_{10} - E(10) = 33 - 10 = 23$$

$$w_{11} = E_{11} - E(11) = 19 - 11 = 8; \quad w_{12} = E_{12} - E(12) = 25 - 12 = 13$$

- Identify the positive number. Using the relation,  
 $u_i = \frac{w_j^2}{v_i}$  for  $i = 1, 2, 3, 4, 5, 6$  and  $j = 7, 8, 9, 10, 11, 12$

$$u_1 = \frac{w_7^2}{v_1} = \frac{8^2}{2} = 8; \quad u_2 = \frac{w_8^2}{v_2} = \frac{13^2}{13} = 13$$

$$u_3 = \frac{w_9^2}{v_3} = \frac{5^2}{5} = 5; \quad u_4 = \frac{w_{10}^2}{v_4} = \frac{23^2}{23} = 23$$

$$u_5 = \frac{w_{11}^2}{v_5} = \frac{8^2}{8} = 8; \quad u_6 = \frac{w_{12}^2}{v_6} = \frac{13^2}{13} = 13$$

- To find the plain text  $A_i = [u_i - \frac{n}{2}] \bmod 26$

$$A_1 = [u_1 - \frac{12}{2}] \bmod 26 = [8 - \frac{12}{2}] \bmod 26 = 2$$

$$A_2 = [u_2 - \frac{12}{2}] \bmod 26 = [13 - \frac{12}{2}] \bmod 26 = 7$$

$$\begin{aligned}
A_3 &= [u_3 - \frac{12}{2}] \bmod 26 = [5 - \frac{12}{2}] \bmod 26 = 25 \\
A_4 &= [u_4 - \frac{12}{2}] \bmod 26 = [23 - \frac{12}{2}] \bmod 26 = 17 \\
A_5 &= [u_5 - \frac{12}{2}] \bmod 26 = [8 - \frac{12}{2}] \bmod 26 = 2 \\
A_6 &= [u_6 - \frac{12}{2}] \bmod 26 = [13 - \frac{12}{2}] \bmod 26 = 7
\end{aligned}$$

Therefore, we deduce from the encoding table that the plain text is **CHURCH**.

## 5. CONCLUSION

In this paper we have identified  $A_n^2, A_n^3, A_n^4$  Arrow graphs and proved that the graphs are RSM labeling. We have applied encryption coding by using a revised Graph Message Jumble Code technique with new labeling and numbering of alphabets based on vowels. In the future, we intend to introduce new labeling technique and prove coding, utilising various graphs in conjunction with various methods of alphabet numbering.

## REFERENCES

- [1] Gallian, J. A, (2010), A Dynamic Survey of Graph labeling, Electron. J. Comb., 18, pp. 56.
- [2] Harary, F., (1988), Graph Theory, Narosa Publishing House Reading, New Delhi.
- [3] Ponraj, R., Somasundaram, S., (2003), Mean Labeling of Graphs, National Academy of Science Letters, 126, pp. 210-213.
- [4] Meena, S., Mani, R., (2019), Root Square Mean Labeling of Some Cycle Related Graphs, IJSART, 5(7), pp. 786-789.
- [5] Meena, S., Mani, R., (2022), Root Square Mean Labeling of theta Related Graphs, MSEA, 71(4), pp. 13049-13070.
- [6] Sandhya, S. S., Somasundaram, S., Anusa, S., (2014), Root Square Mean Labeling of Graphs, Int. J. Contemp. Math. Sci., 9(14), pp. 667-676.
- [7] Sandhya, S. S., Somasundaram, S., Anusa, S., (2014), Root Square Mean Labeling of Some New Disconnected Graphs, IJMTT, 15(2), pp. 85-92.
- [8] Caminiti, S., Finocchi, I., Peterschi, R., (2007), On Coding Labeled Trees, Theoretical Computer Science, 382, pp. 97-108.
- [9] Uma Maheswari, A., Ambika, C., (2022), New coding Algorithms using AUM block sum labeling, Neuroquantology, 20(9), pp. 377-385.



**Prof. T. Christy** is an Assistant Professor in the Department of Mathematics at Patrician College of Arts and Science in Chennai, Tamil Nadu, India. She has teaching experience at various institutions. She is skilled in developing new lessons, providing conceptual knowledge, and has strong analytical and mathematical abilities.



**Dr. G. Palani** completed his research work from Anna University, Chennai, India in 2001. He received BK21 fellowship from Inha University, South Korea. Presently, he is working as Associate Professor, Department of Mathematics, Dr Ambedkar Govt Arts College, Chennai-39, India. He published more than 85 papers in International/ National Reputed Journals. His Area of research work is Applied Mathematics and Graph Theory.